

Finding an integer pair with given sum is a very familiar problem in algorithm community. The general problem is: given an array of integer numbers and a given sum, how do we find out a integer pair with given sum?

The easiest solution is to consider every possible pair and compare it with the given number. Time complexity of this approach will be $O(n^2)$ which is too big if the array size is very big.

Most efficient solution is to applying an efficient sorting algorithm to sort the numbers first. You can use [QuickSort](#) which has average case [time complexity](#) of $O(n \log n)$. Once the list is sorted, we can use a linear search on the array to find out the pair with a given sum. Use the following pseudo code.

1. initialize two pointer: **left** & **right** pointing to the first and last element of the sorted array
2. if $\text{array}[\text{left}] + \text{array}[\text{right}] == \text{SUM}$ then return **true**
3. else if $\text{array}[\text{left}] + \text{array}[\text{right}] < \text{SUM}$ then **left++**
4. else **right--**
5. if **left** < **right** then go to step 2
6. else return "**no pair found**"

Why this algorithm works? The intuition is simple. The main technique is to move the left and right pointer. If $\text{array}[\text{left}] + \text{array}[\text{right}] < \text{SUM}$ that means we have to take a bigger number from the array to form the given sum. So we move the left pointer in right direction to form a bigger sum. As the array is sorted in increasing order, by moving the left pointer we'll get a bigger number and hence bigger sum. On the other hand if $\text{array}[\text{left}] + \text{array}[\text{right}] > \text{SUM}$, then we'll have to take a smaller number to form the given sum. We do that by moving the right pointer to the left where we'll get a smaller number.